

Adaptive Parameterized Consistency for Non-Binary CSPs by Counting Supports

R.J.Woodward^{1,2}, A.Schneider¹, B.Y.Choueiry¹, and C.Bessiere²

¹Constraint Systems Laboratory • University of Nebraska-Lincoln • USA

²LIRMM-CNRS • University of Montpellier • France

1. Contributions

1. Extend p -stability for AC to GAC (to operate on non-binary CSPs)
2. Provide a more precise mechanism for computing it
3. Algorithm for enforcing apc -LC

2. Local Consistency

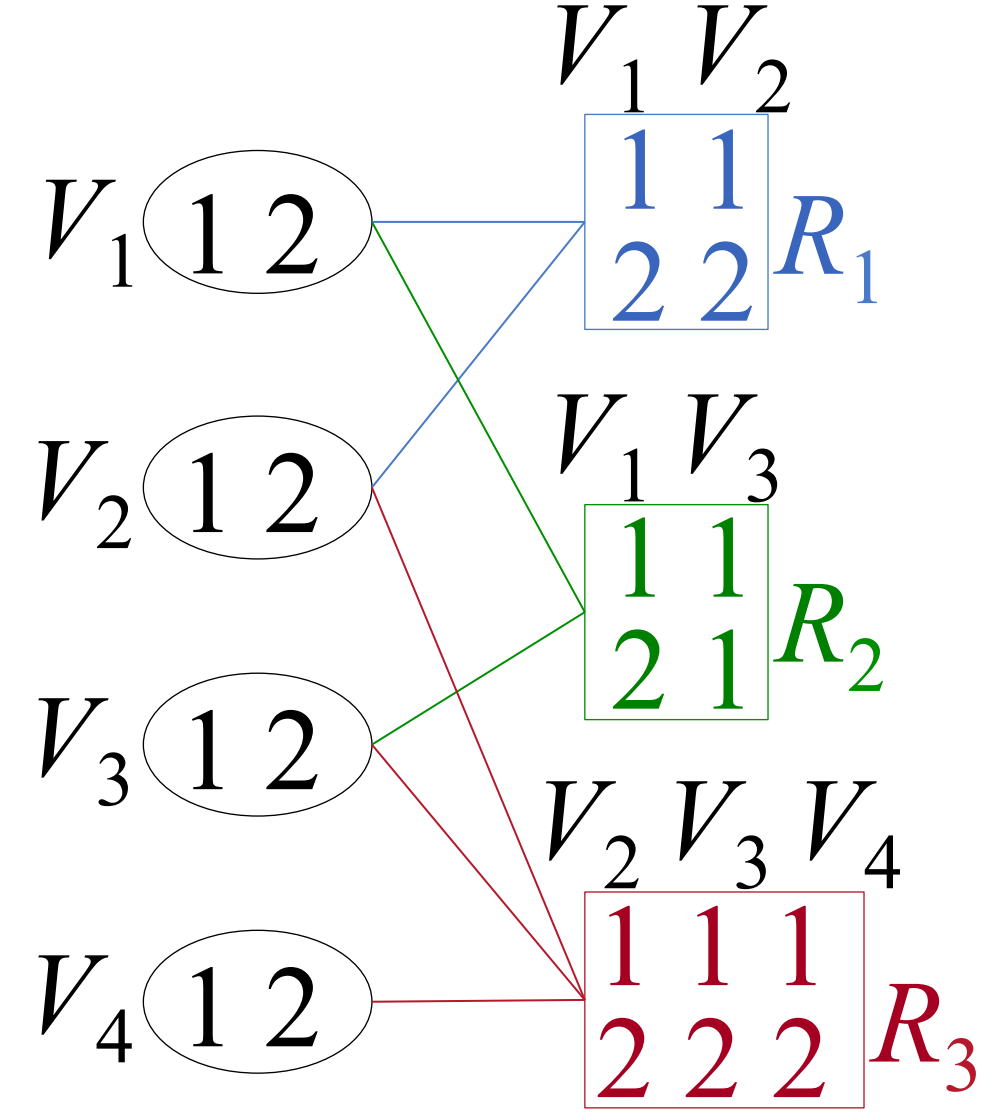
Generalized Arc Consistency (GAC)

ensures any value in the domain of any variable in the scope of every relation can be extended to a tuple satisfying the relation. E.g., filters value 2 from V_3 .

STR ensures GAC, filtering both domains and relations. E.g., filters value 2 from V_3 , tuple $\langle 2,1 \rangle$ from R_2 .

m -wise Consistency, $R^{*}(m)C$ ensures every tuple in every relation can be extended to every combination of $m-1$ other relations. $R^{*}(3)C$ filters tuple $\langle 2,2 \rangle$ from R_1 , $\langle 2,1 \rangle$ from R_2 and $\langle 2,2,2 \rangle$ from R_3 . Projecting relations onto domains filters 2 from every variable.

Pairwise Consistency $\equiv R^{*}(2)C$.



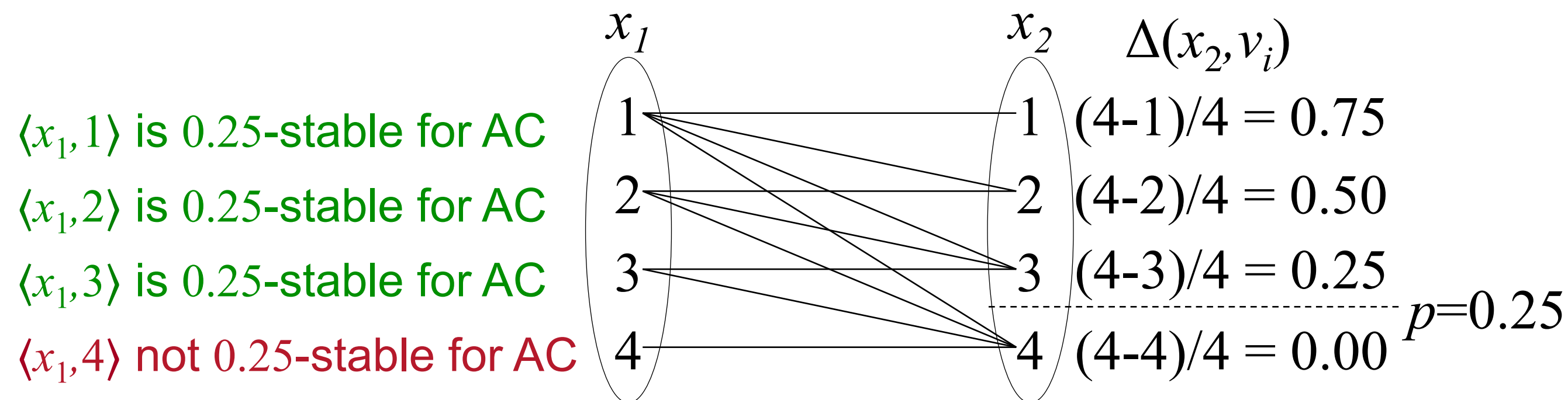
3. Parameterized Adaptive Consistency

p -stability for AC [Balafrej+ CP2013]

Relative position of v_i in $dom^o(x_i)$:

$$\Delta(x_i, v_i) = \frac{|dom^o(x_i)| - rank(v_i, dom^o(x_i))}{|dom^o(x_i)|}$$

For parameter p , $\langle x_i, v_i \rangle$ is p -stable for AC for c_{ij} on x_i, x_j if it has AC-support $\langle x_i, v_j \rangle$ with $\Delta(x_i, v_j) \geq p$.



Issue: Position is not a precise measure of support.

p -stability for GAC

$\langle x_i, v_i \rangle$ is p -stable for GAC if for every constraint c_j on x_i

$$\frac{|\sigma_{x_i=v_i}(R_j)|}{|R_j^o|} \geq p$$

Where R_j^o is the original, unfiltered relation.

$gacSupports[R_j](\langle x_i, v_i \rangle)$: data structure to store $\sigma_{x_i=v_i}(R_j)$.

$gacSupports[R_1](\langle x_1, 1 \rangle) = \{0, 1, 2, 3\}$ ($x_1, 1$ is 0.25-stable for GAC ($\frac{4}{10} \geq 0.25$))
 $gacSupports[R_1](\langle x_1, 2 \rangle) = \{4, 5, 6\}$ ($x_1, 2$ is 0.25-stable for GAC ($\frac{3}{10} \geq 0.25$))
 $gacSupports[R_1](\langle x_1, 3 \rangle) = \{7, 8\}$ ($x_1, 3$ not 0.25-stable for GAC ($\frac{2}{10} < 0.25$))
 $gacSupports[R_1](\langle x_1, 4 \rangle) = \{9\}$ ($x_1, 4$ not 0.25-stable for GAC ($\frac{1}{10} < 0.25$))

$x_j \backslash x_i$	1	2	3	4
0	1	1		
1	1	2		
2	1	3		
3	1	4		
4	2	2		
5	2	3		
6	2	4		
7	3	3		
8	3	4		
9	4	4		

Parameterized Adaptive Consistency

- LC: A local consistency stronger than (G)AC
- pc -LC: $\forall \langle x_i, v_i \rangle$ and c_j on x_i , $\langle x_i, v_i \rangle$ is
 - p -stable for (G)AC on c_j or
 - LC on c_j
- apc -LC: adaptive version of pc -LC that uses the weight of a constraint, $w(c_j)$, like $dom/wdeg$

$$p(c_j) = \frac{w(c_j) - \min_{c_k \in C}(w(c_k))}{\max_{c_k \in C}(w(c_k)) - \min_{c_k \in C}(w(c_k)) + 1}$$

4. Algorithm & Empirical Evaluations

- Generate the $gacSupports$ at pre-processing
- $\forall c_j, \forall x_i$ in $c_j, \forall v_i$ in $dom(x_i)$
 - Enforce LC when $\langle x_i, v_i \rangle$ is not p -stable for GAC
 - Enforce STR when $\langle x_i, v_i \rangle$ has no support

Algorithm 1: LIVING-STR(c_i): set of variables

Input: c_j : a constraint of \mathcal{P}
Output: Set of variables in $scope(c_j)$ whose domains have been modified

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1  $X_{modified} \leftarrow \emptyset$ 
2 foreach  $x_i \in scope(c_j) \mid x_i \notin past(\mathcal{P})$  do
3   foreach  $v_i \in dom(x_i)$  do
4     if  $|gacSupports[R_j](\langle x_i, v_i \rangle)| \neq 0$  and  $\frac{|gacSupports[R_j](\langle x_i, v_i \rangle)|}{|R_j^o|} \not\geq p(c_j)$  then
5       APPLY-LC( $R_j, gacSupports[R_j](\langle x_i, v_i \rangle)$ )
6     if  $|gacSupports[R_j](\langle x_i, v_i \rangle)| = 0$  then
7       foreach  $c_k \in cons(x_i)$  do
8         delTuples( $c_k, gacSupports[R_k](\langle x_i, v_i \rangle), |past(\mathcal{P})|$ )
9          $dom(x_i) \leftarrow dom(x_i) \setminus \{v_i\}$ 
10        if  $dom(x_i) = \emptyset$  then throw INCONSISTENCY
11         $X_{modified} \leftarrow X_{modified} \cup \{x_i\}$ 
12 return  $X_{modified}$ 
    
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Experimental Results

apc - $R^{*}(2)C$ can solve more instances and quicker than both STR and $R^{*}(2)C$.

Total number of instance: 623	STR	$R^{*}(2)C$	apc - $R^{*}(2)C$
Number of instances solved by	504	550	552
# instances solved only by	10	5	0
# instances solved by STR, but missed by	0	18	11
# instances solved by $R^{*}(2)C$, but missed by	64	0	6
# instances solved by apc - $R^{*}(2)C$, but missed by	59	8	0
Average CPU time (sec.) over 458 instances	328.41	378.12	313.31
Median CPU time (sec.) over 458 instances	7.23	17.35	7.21

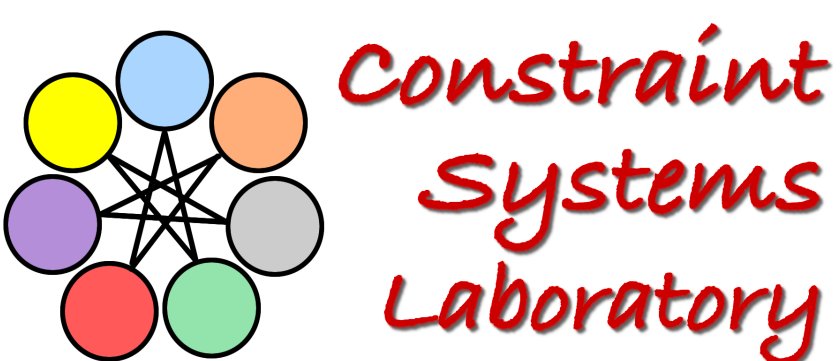
Benchmark results where apc - $R^{*}(2)C$ performs the best, is competitive, and is the worst. Also included, a benchmark that could not be solved by STR.

benchmark	# Instances	Completed			Average CPU Time (sec)			Median CPU Time (sec)		
		STR	$R^{*}(2)C$	apc - $R^{*}(2)C$	STR	$R^{*}(2)C$	apc - $R^{*}(2)C$	STR	$R^{*}(2)C$	apc - $R^{*}(2)C$
a) apc - $R^{*}(2)C$ is the best										
aim-50	24	24	24	24	0.04	0.07	0.04	0.02	0.04	0.03
allIntervalSeries	25	22	22	22	7.09	141.85	6.00	0.13	0.31	0.12
inhSat	16	16	16	16	13.07	357.66	11.74	8.15	142.24	7.21
modifiedRenault	50	50	50	50	6.39	11.17	6.29	7.24	8.79	6.98
rand-3-20-20	50	31	43	41	1,666.10	939.88	932.77	1,211.50	822.54	811.74
b) apc - $R^{*}(2)C$ is competitive										
aim-100	24	24	24	24	0.38	0.26	0.41	0.18	0.25	0.16
aim-200	24	22	24	22	414.48	6.52	286.27	2.39	1.37	2.60
inhUnsat	34	34	34	34	13.61	294.77	13.95	10.74	153.50	9.78
lexVg	63	63	63	63	69.81	341.87	338.74	0.50	1.38	0.89
pret	8	4	4	4	117.89	347.03	136.04	115.81	354.82	145.70
rand-3-20-20-fed	50	39	48	47	928.06	546.84	615.23	501.30	422.24	464.00
rand-8-20-5	20	9	20	20	2,564.94	355.57	372.76	1,987.35	314.26	261.68
rand-10-20-10	20	12	12	12	6.72	1.67	2.76	6.40	1.66	2.75
ssa	8	6	5	6	64.60	100.64	69.59	1.51	1.60	1.58
TSP-25	15	13	10	13	232.38	1,072.72	743.33	69.00	211.41	131.69
ukVg	65	37	31	34	166.82	796.90	421.35	36.29	54.65	30.39
varDimacs	9	6	6	6	89.23	587.55	319.20	1.56	6.43	2.94
wordsVg	65	65	58	58	119.76	532.05	400.22	0.39	0.95	0.59
c) apc - $R^{*}(2)C$ is the worst										
dubois	13	7	8	6	1,000.54	451.91	1,456.01	552.13	255.25	779.57
TSP-20	15	15	15	15	101.20	318.37	335.13	23.32	61.55	46.34
d) Not solved by STR										
dag-rand	25	0	25	25	-	123.70	149.64	-	124.47	151.33

Average number of calls to STR and $R^{*}(2)C$ in apc - $R^{*}(2)C$

- On *allIntervalSeries*, no $R^{*}(2)C$ calls because solved at pre-processing (no weight set).
- Sometimes, there are more STR calls, others, more $R^{*}(2)C$ calls. Thus, apc - $R^{*}(2)C$ finds a good mix between STR and $R^{*}(2)C$ for each instance. Success!

benchmark	STR Calls	$R^{*}(2)C$ Calls	benchmark	STR Calls	$R^{*}(2)C$ Calls
a) apc - $R^{*}(2)C$ is the best					
aim-50	456,823	39,491	aim-100	7,731,585	894,353
allIntervalSeries	38,281,694	0	aim-200	1,160,334,482	163,177,907
inhSat	22,119,135	599,080	inhUnsat	51,688,166	1,918,781
modifiedRenault	4,618,778	601,641	lexVg	564,010,457	2,180,503,026
rand-3-20-20	489,441,126	3,480,216,943	pret	422,987,946	13,973,748
b) apc - $R^{*}(2)C$ is competitive					
rand-3-20-20-fed	455,664,100	2,956,467,994	rand-3-20-20-fed	455,664,100	2,956,467,994
c) apc - $R^{*}(2)C$ is the worst					
dubois	3,343,830,604	4,668,288	rand-8-20-5	77,470,561	184,764,543
TSP-20	622,949,698	991,590,957	rand-10-20-10	72,608	3,972
d) Not solved by STR					
dag-rand	359,248	21,870	ssa	156,631,370	11,689,961
			TSP-25	2,903,953,315	3,947,391,769
			ukVg	341,565,892	1,002,334,753
			varDimacs	720,843,958	84,123,204
			wordsVg	514,840,737	2,052,367,934



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